# Reducer-Based Construction of Conditional Verifiers

#### **Dirk Beyer**

#### LMU Munich, Germany

#### @ CPAchecker/LDV Workshop, Moscow, 2018-09-25







### Many Verification Tools Available



### Vision

#### I have a dream ...

- ... that one day, all tools for formal methods work together to solve hard verification problems and make our world safer and more secure.
- In that one day, model checkers and theorem provers can be integrated into the software-development process as seamless as unit testing today.
- … that one day, model checkers, theorem provers, SMT solvers, and testers use common interfaces for interaction and composition.

### Outline

Dream is not utopian — there are a few approaches already ...

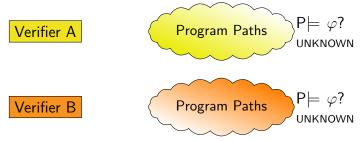
- Approach 1: Conditional Model Checking [FSE'12]
- Approach 2: Verification Witnesses [FSE'15, FSE'16]
- Approach 3: Tests from Witnesses [TAP'18]

...

Cooperative Verification by Conditional Model Checking and Reducers

### Facing Hard Verification Tasks

Given: Program  $\mathsf{P}\models\varphi$ ?



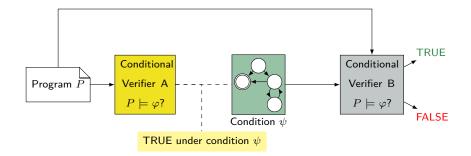
# Facing Hard Verification Tasks

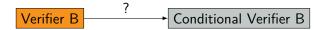
Given: Program  $P \models \varphi$ ?  $\mathsf{P}\models\varphi?$ Verifier A **Program** Paths UNKNOWN  $\mathsf{P} \models \varphi?$ Verifier B **Program Paths** UNKNOWN Verifier A + Verifier B **Program Paths**  $\mathsf{P}\models\varphi\checkmark$ 

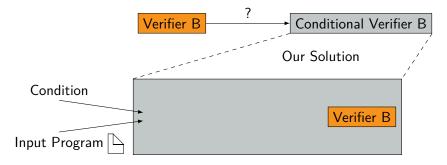
e.g., conditional model checking

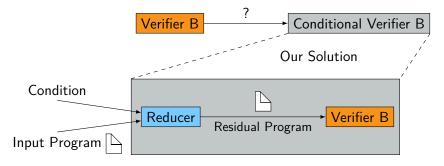
# Conditional Model Checking

[Beyer/Henzinger/Keremoglu/Wendler FSE'12, DOI Link, Preprint Link]]



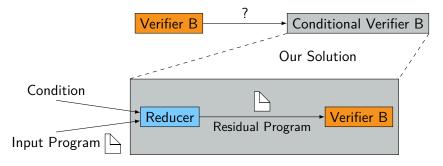






#### Reducer (preprocessor)

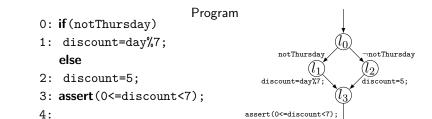
- Builds standard input (C program)
- Representing a subset of paths
- Contains at least all non-verified paths



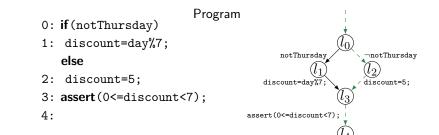
#### Reducer (preprocessor)

- Builds standard input (C program)
- Representing a subset of paths
- Contains at least all non-verified paths
- + Verifier-unspecific approach
- + Many conditional verifiers possible

# Example Program and Condition

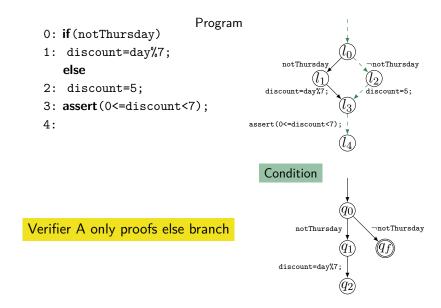


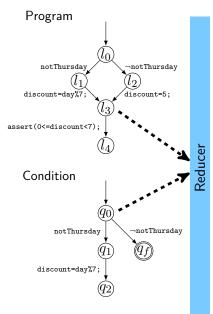
# Example Program and Condition



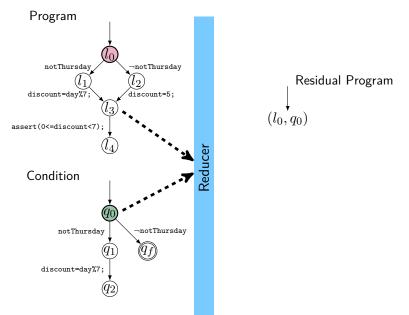
#### Verifier A only proofs else branch

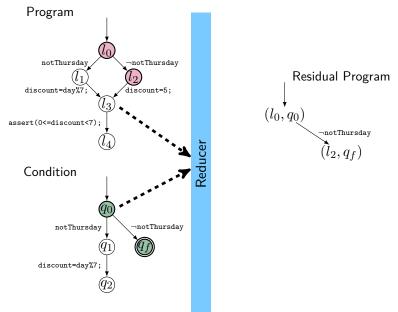
# Example Program and Condition

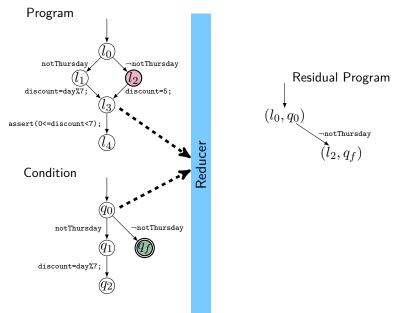


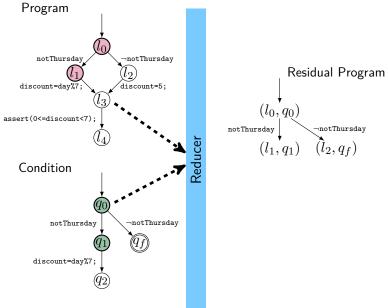


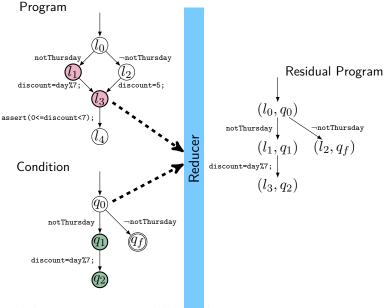
**Residual Program** 

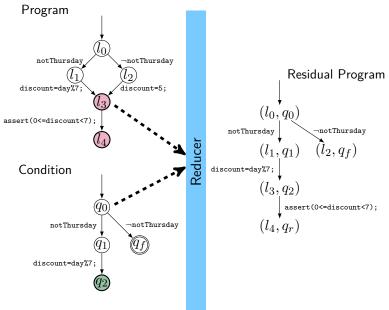




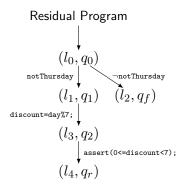




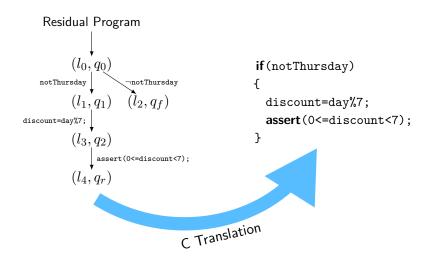




#### Reducer: C Transformation

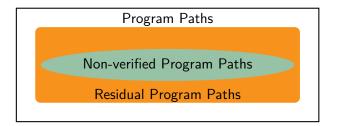


#### Reducer: C Transformation



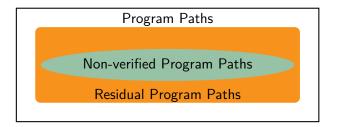
### Reducer: Soundness

#### **Residual Condition**



### Reducer: Soundness

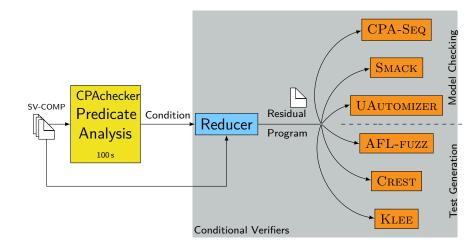
#### **Residual Condition**



#### Theorem

Presented reducer fulfills residual condition.

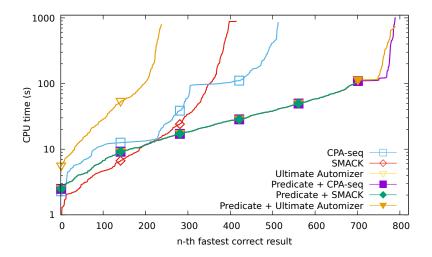
### **Evaluation Setup**



### Small Extract of Results

|    |                             |  |  | Pre   | DICATE   | Pree   | DICATE   |  |
|----|-----------------------------|--|--|---|--|--|--|--|
| CP | CPA-Seq                     |  | UAUTOMIZER   |   | +Reducer   |  | +Reducer   |  |
|    |                             |  |  | +c  | PA-Seq   | +UA  | UTOMIZER   |  |
| S  | t(s)                        | S  | t(s)   | S   | t(s)   | S  | t(s)   |  |
| X  | 910                         | X  | 900  | <ul> <li>Image: A start of the start of</li></ul> | 120  | 1  | 130  |  |
| X  | 910                         | X  | 910  | 1   | 450  | X  | 1100   |  |
| X  | 950                         | X  | 490  | X   | 910  | 1  | 260  |  |
| X  | 950                         | X  | 910  | X   | 1100   | 1  | 470  |  |
|    | CP<br>S<br>X<br>X<br>X<br>X | S         t(s)           X         910           X         910           X         950 | S       t(s)       S         X       910       X         X       910       X         X       950       X | S       t(s)       S       t(s)         X       910       X       900         X       910       X       910         X       950       X       490   | CPA-SEQ       UAUTOMIZER       +R         -       -       +C         S       t(s)       S       t(s)         X       910       X       900       ✓         X       910       X       910       ✓         X       950       X       490       X | $ \begin{array}{ c c c c c c c } & & & & & & & + \mathrm{CPA-SeQ} \\ \hline S & t(s) & S & t(s) & S & t(s) \\ \hline X & 910 & X & 900 & \checkmark & 120 \\ \hline X & 910 & X & 910 & \checkmark & 450 \\ \hline X & 950 & X & 490 & X & 910 \\ \hline \end{array} $ | $\begin{array}{ c c c c c c c c c c c c c c c c c c c$ |  |

### Effectiveness on Hard Tasks



# More Information: Reducer-Based Construction of Conditional Verifiers

[Proc. ICSE 2018, pages 1182–1193, ACM. DOI Link, Preprint Link]

#### Dirk Beyer, Marie-Christine Jakobs, Thomas Lemberger, and Heike Wehrheim

#### LMU Munich, Germany and Paderborn University, Germany

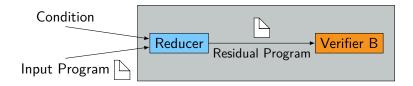




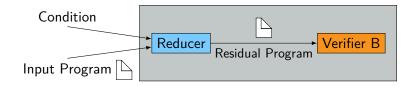




Template-based conditional verifier construction

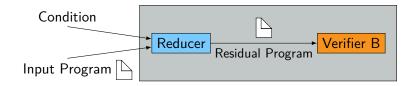


Template-based conditional verifier construction

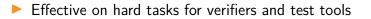


- One Reducer
  - Proven sound
  - Used in many conditional verifiers

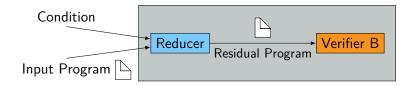
Template-based conditional verifier construction



- One Reducer
  - Proven sound
  - Used in many conditional verifiers



Template-based conditional verifier construction



- One Reducer
  - Proven sound
  - Used in many conditional verifiers
- Effective on hard tasks for verifiers and test tools

#### Future Work

- More reducers
- Using conditions from other tools

Dirk Beyer

LMU Munich, Germany

# Overview Approaches for Combinations

