Assume but Verify: Deductive Verification of Leaked Information in Concurrent Applications Toby Murray, Mukesh Tiwari, Gidon Ernst and David A. Naumann

ACM CCS 2023, November 28, 2023







LUDWIG-MAXIMILIAN



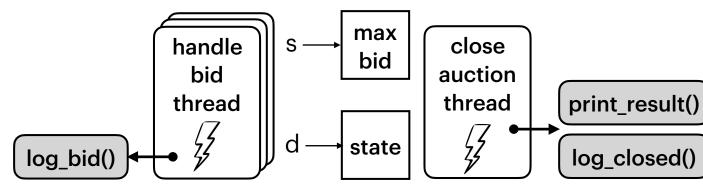


tl;dr

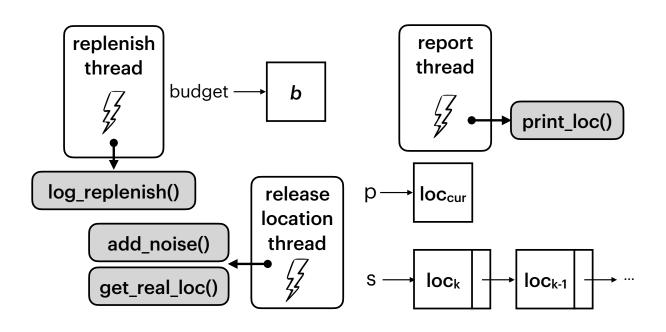
Auto-Active Verification for Concurrent C Programs against Declarative Dynamic Declassification Security Policies via Security Concurrent Separation Logic reasoning against a Constant Time threat model

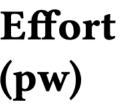
Case Study	Proof Ratio	Verified SLOC	Unverified SLOC	E (j
Location Service	1.9	210	124	3
Auction Server	4.3	187	79	3
Wordle	5.9	47	81	0
Private Learning	2.5	315	114	6

Sealed Bid Auction Server



Private Location Service





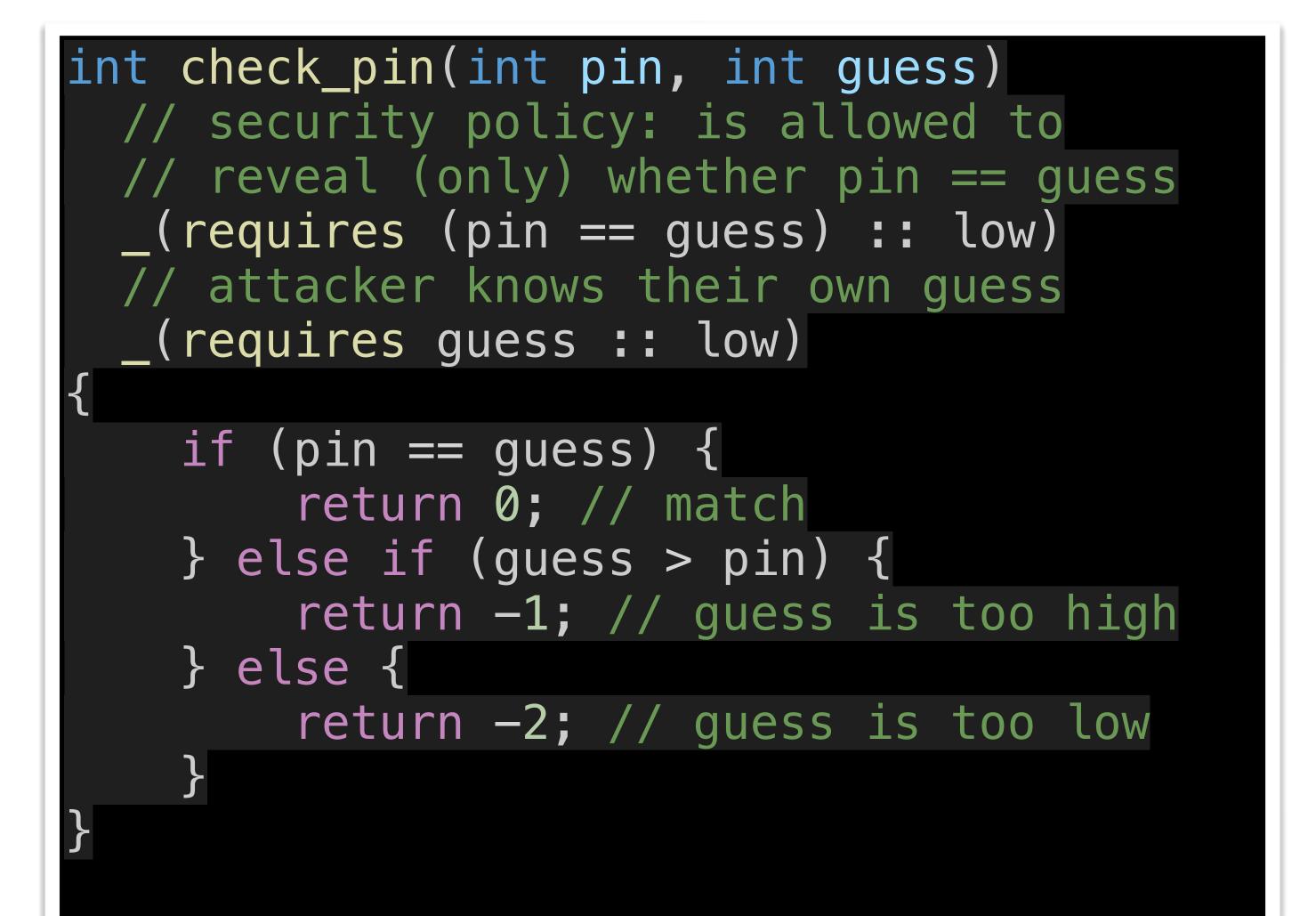


- 3.5
- 0.3
- 6



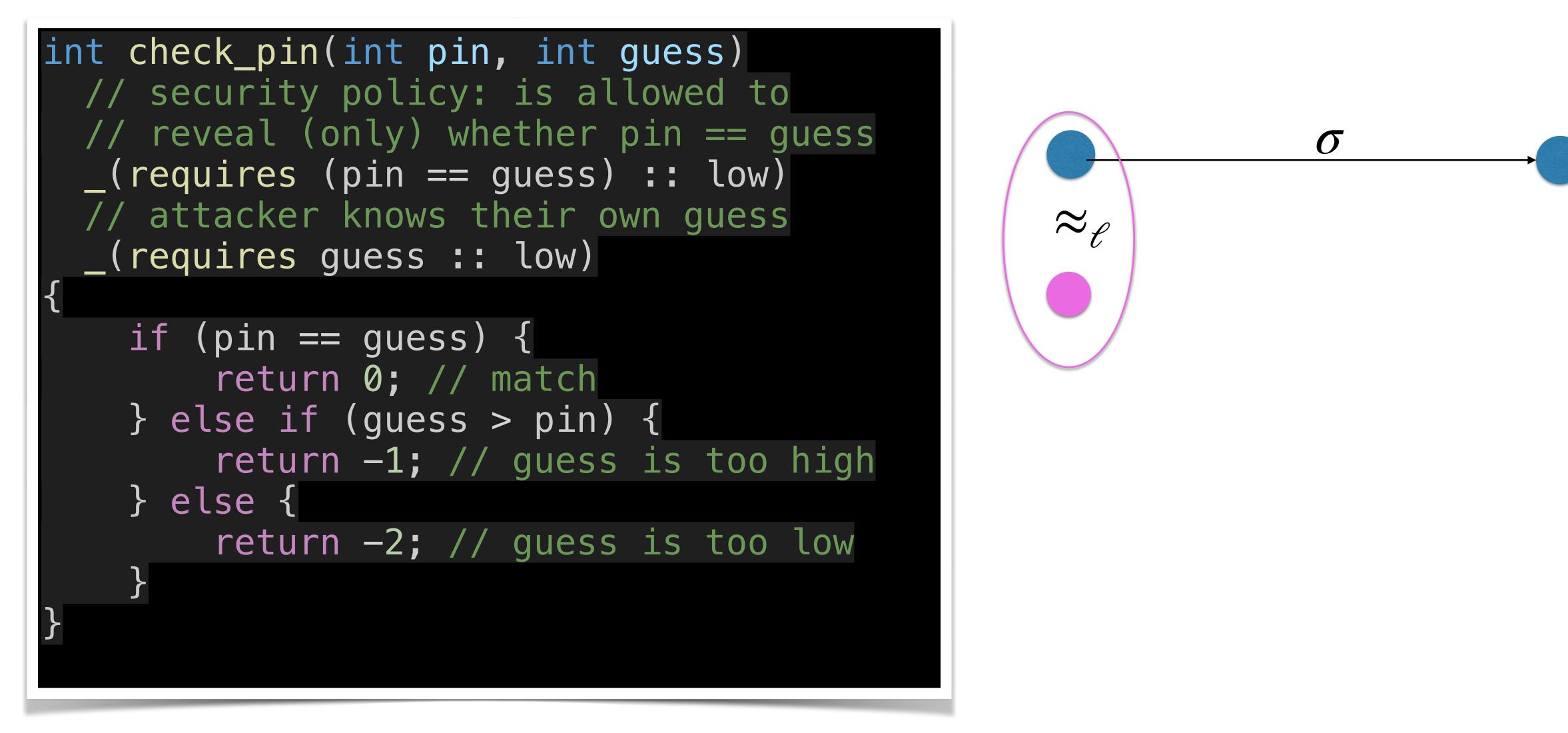
Private Learning



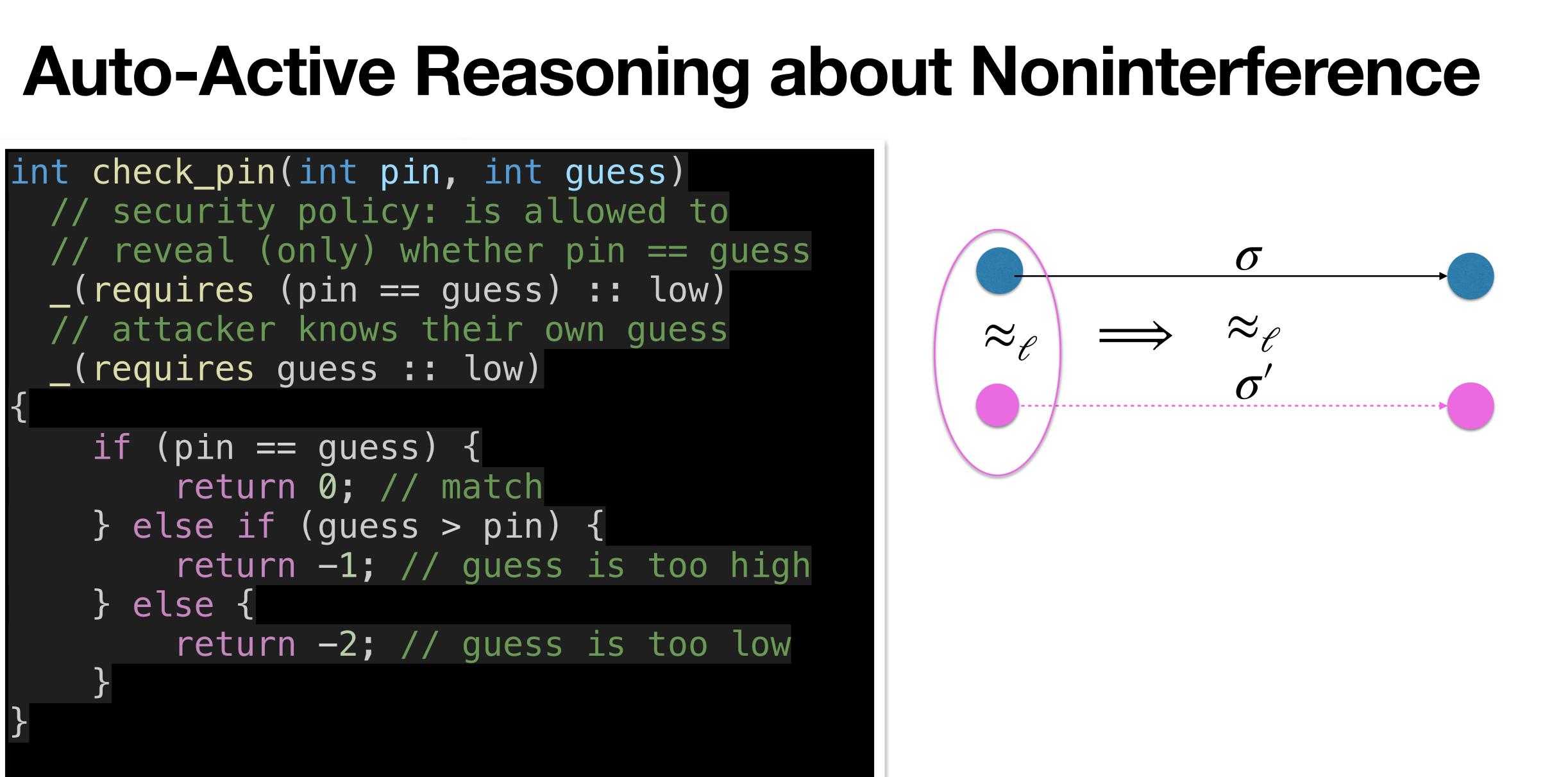


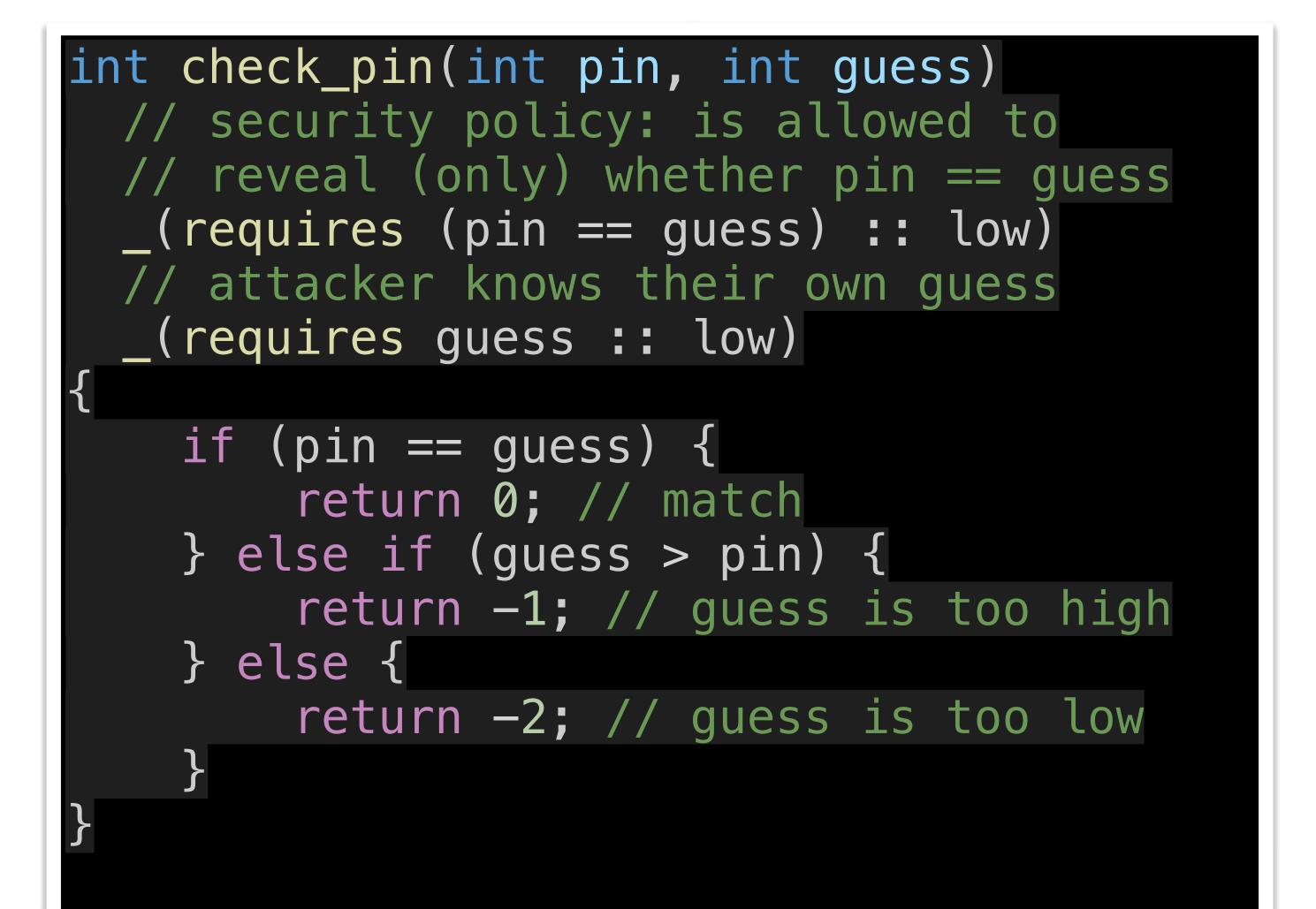


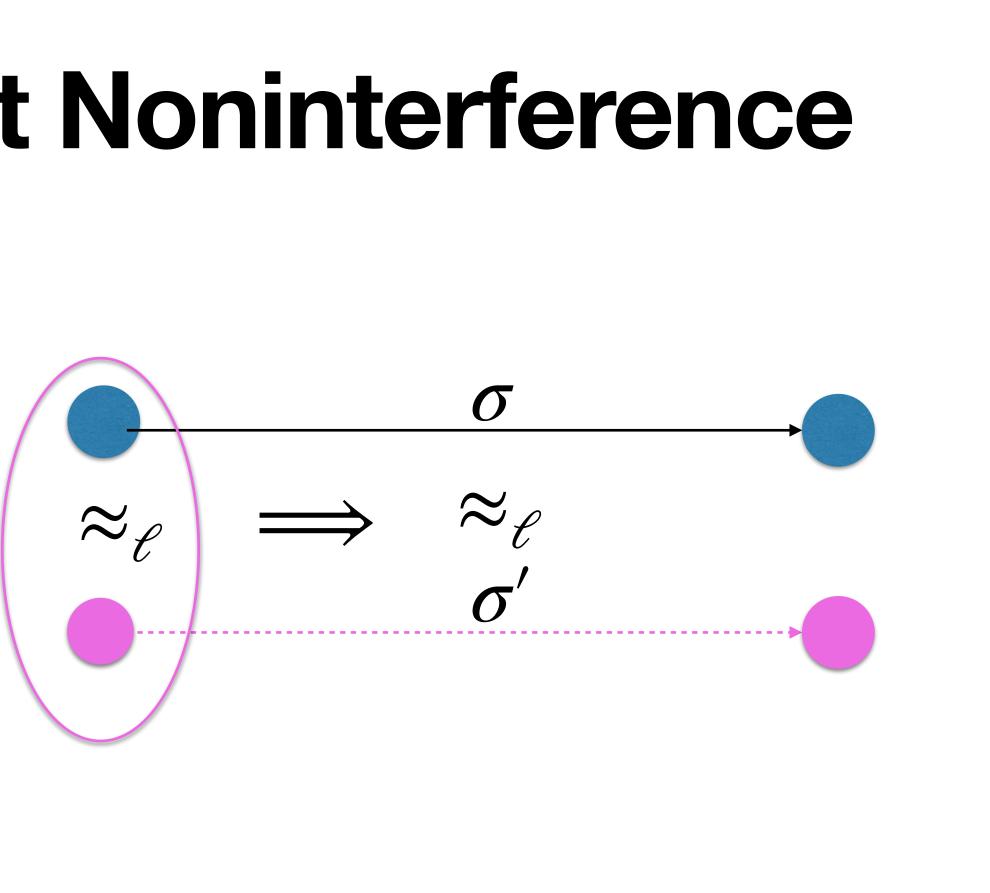




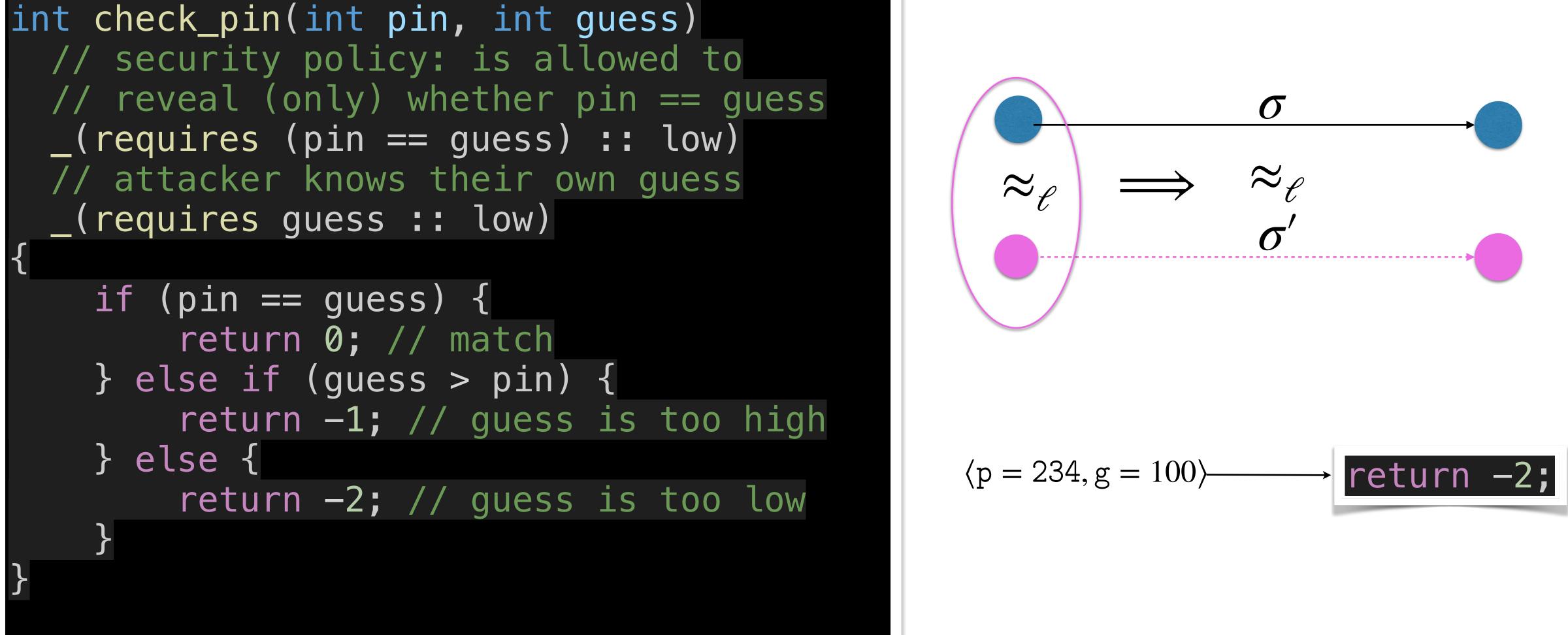


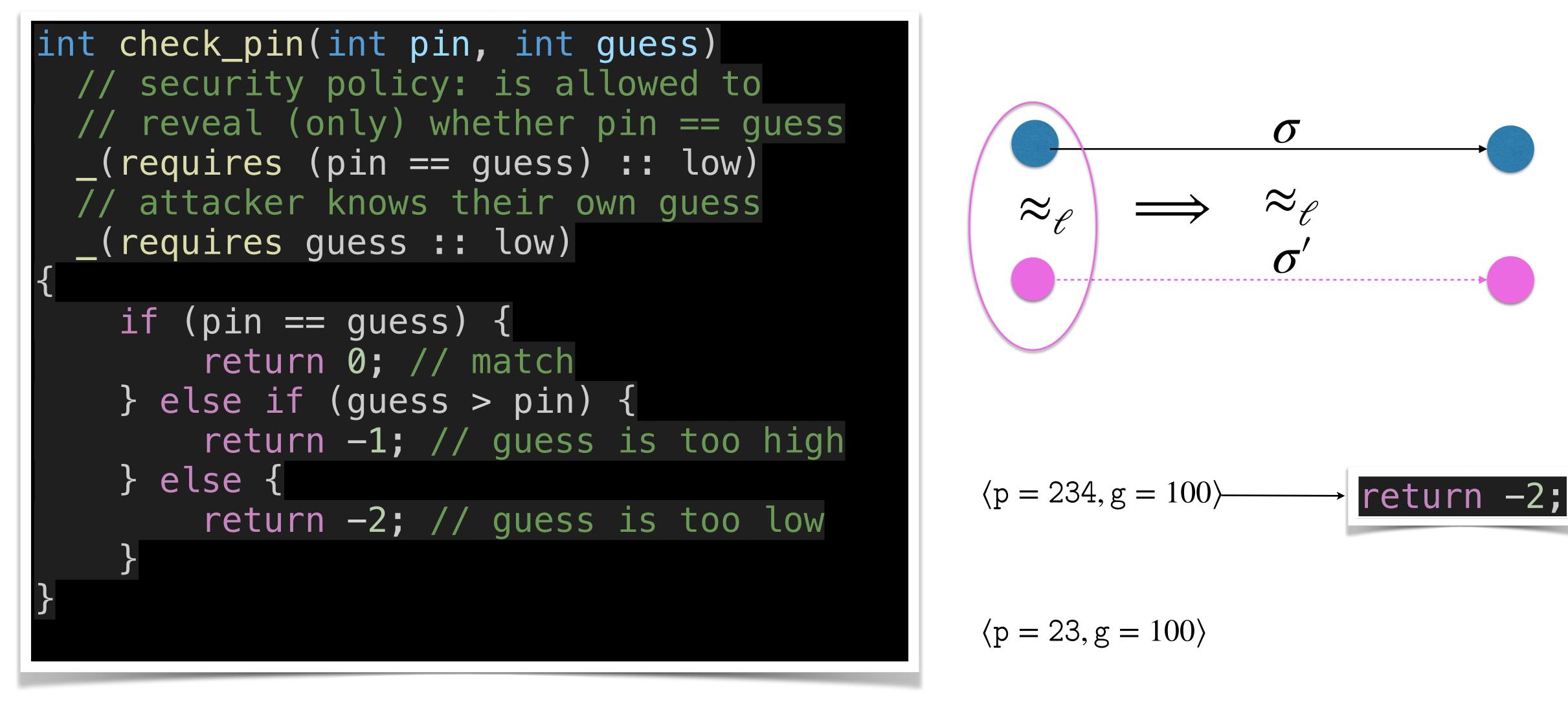




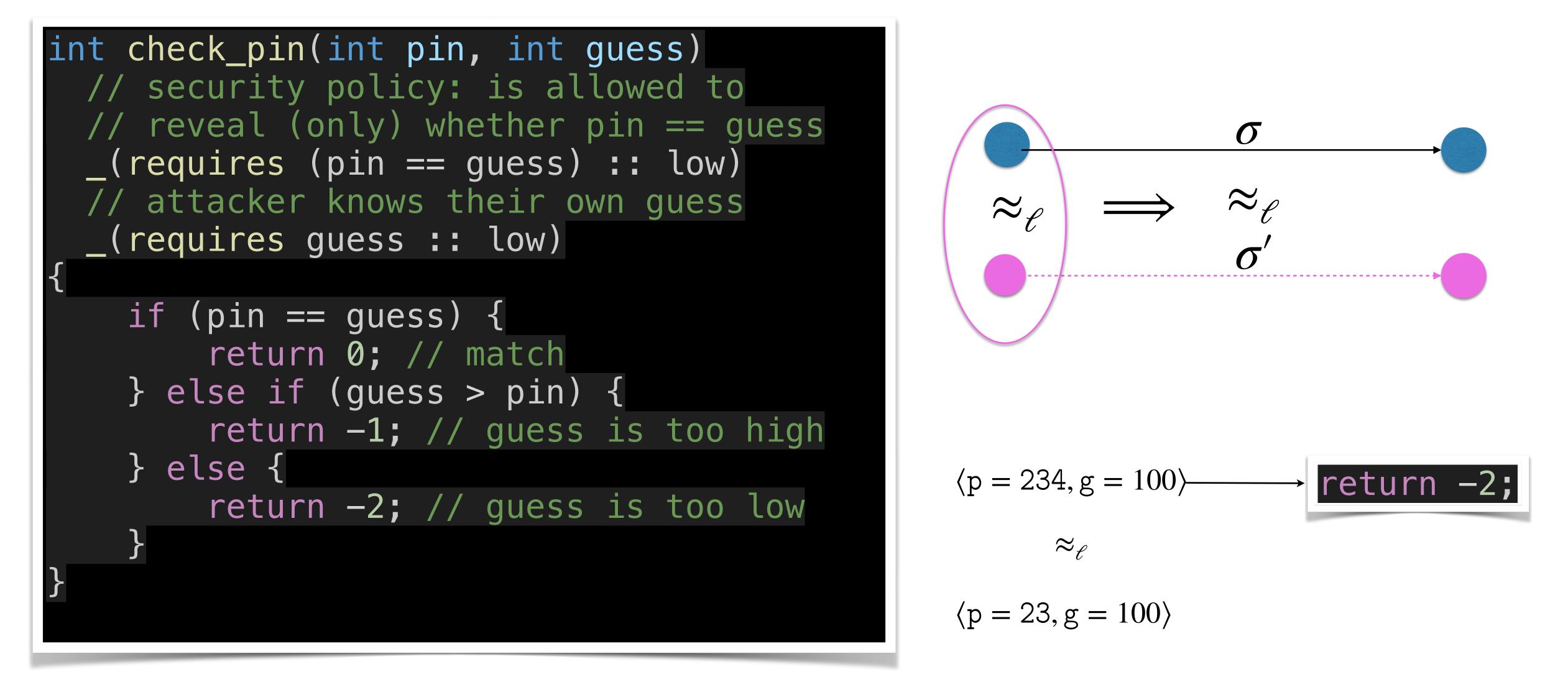


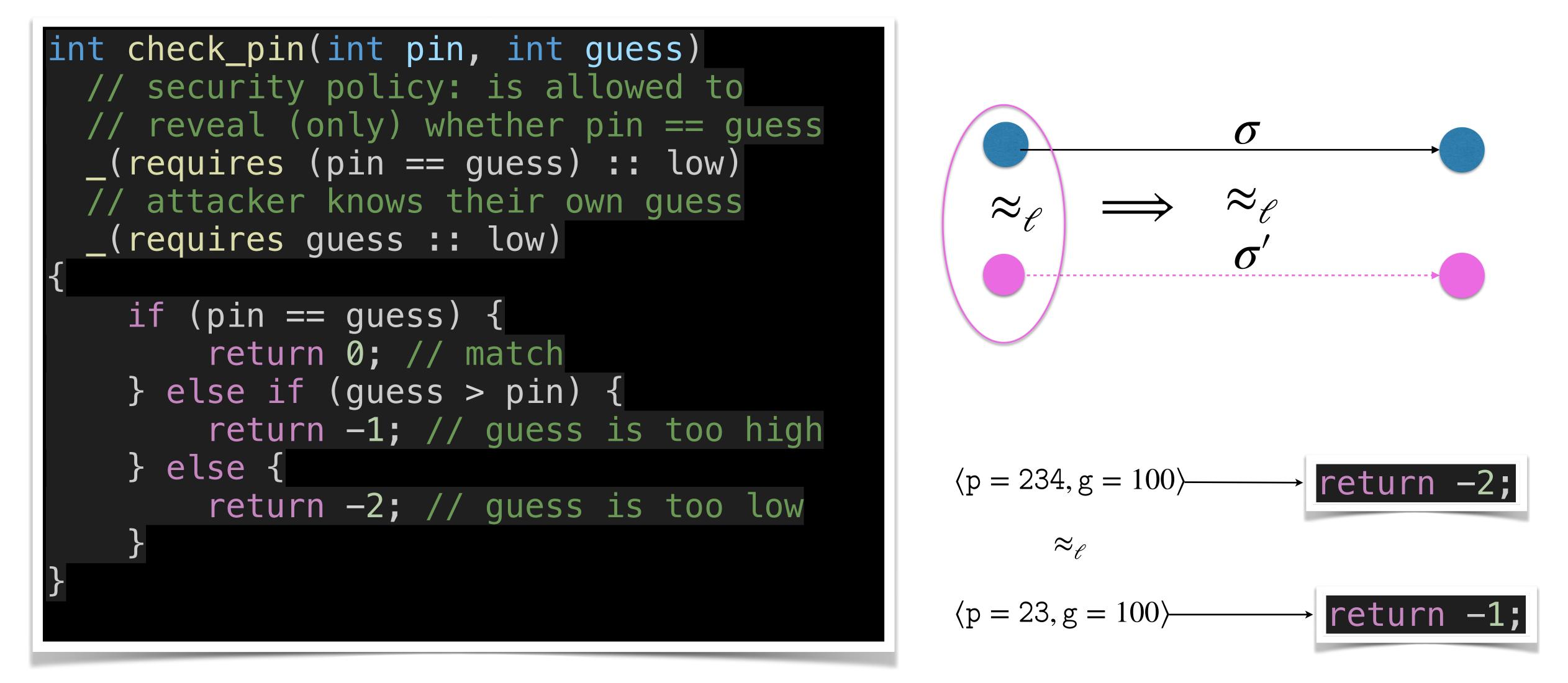
 $\langle p = 234, g = 100 \rangle$

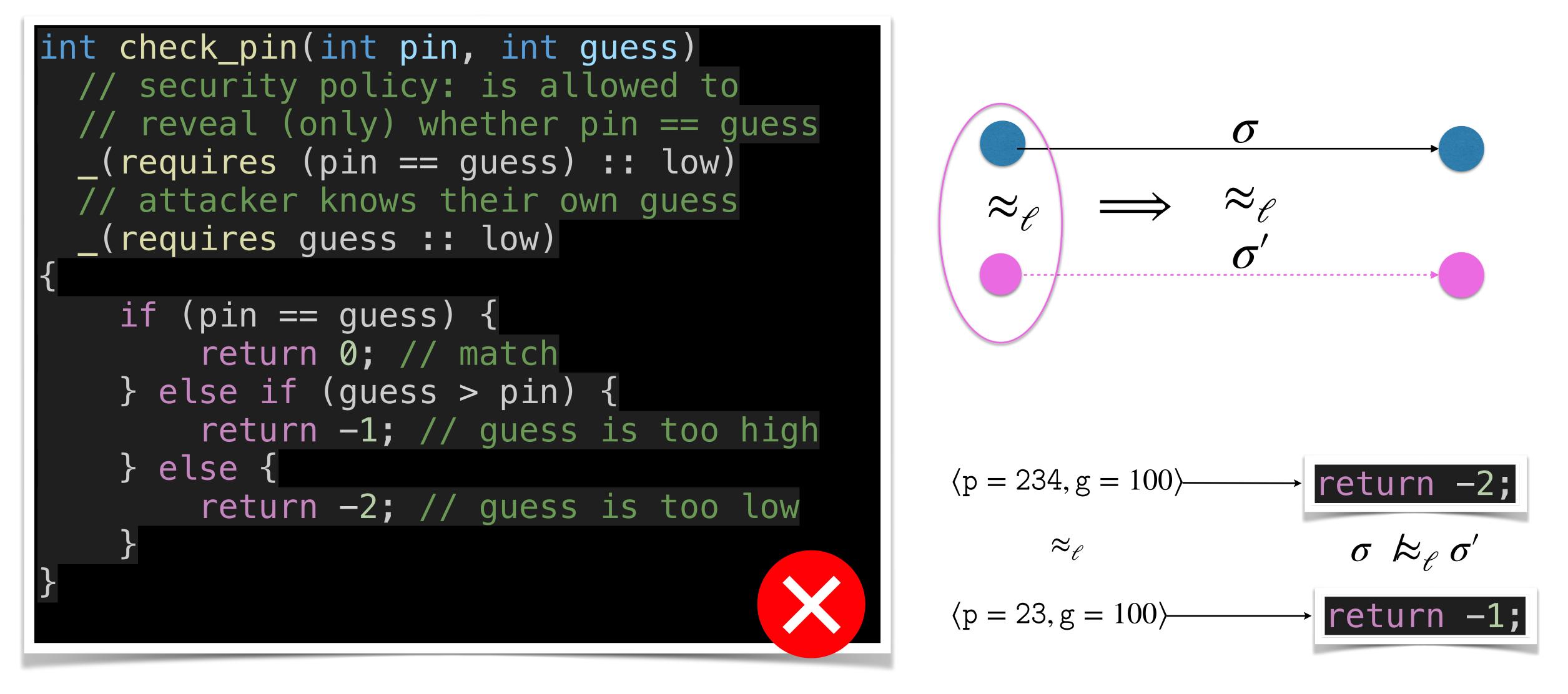




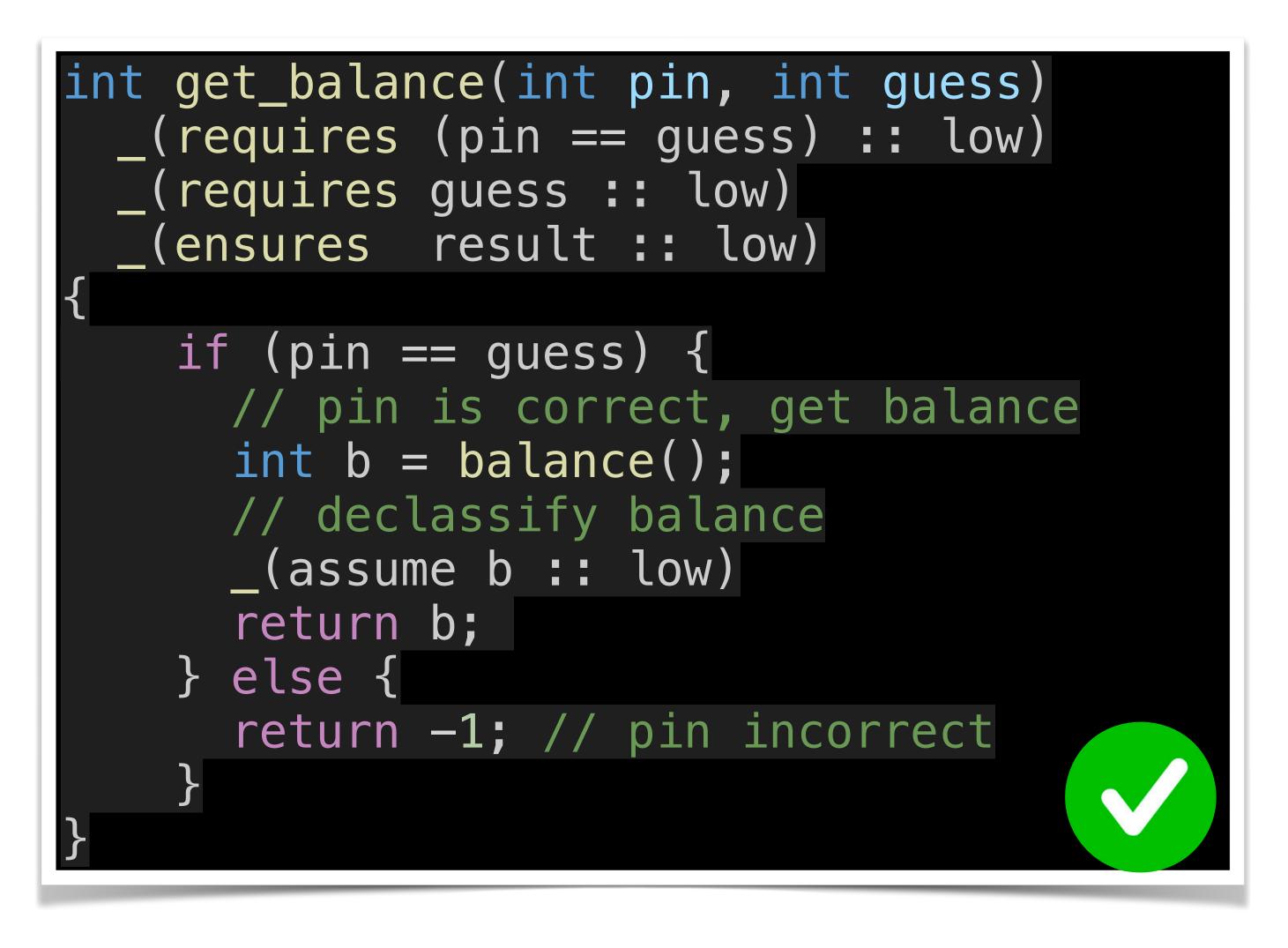






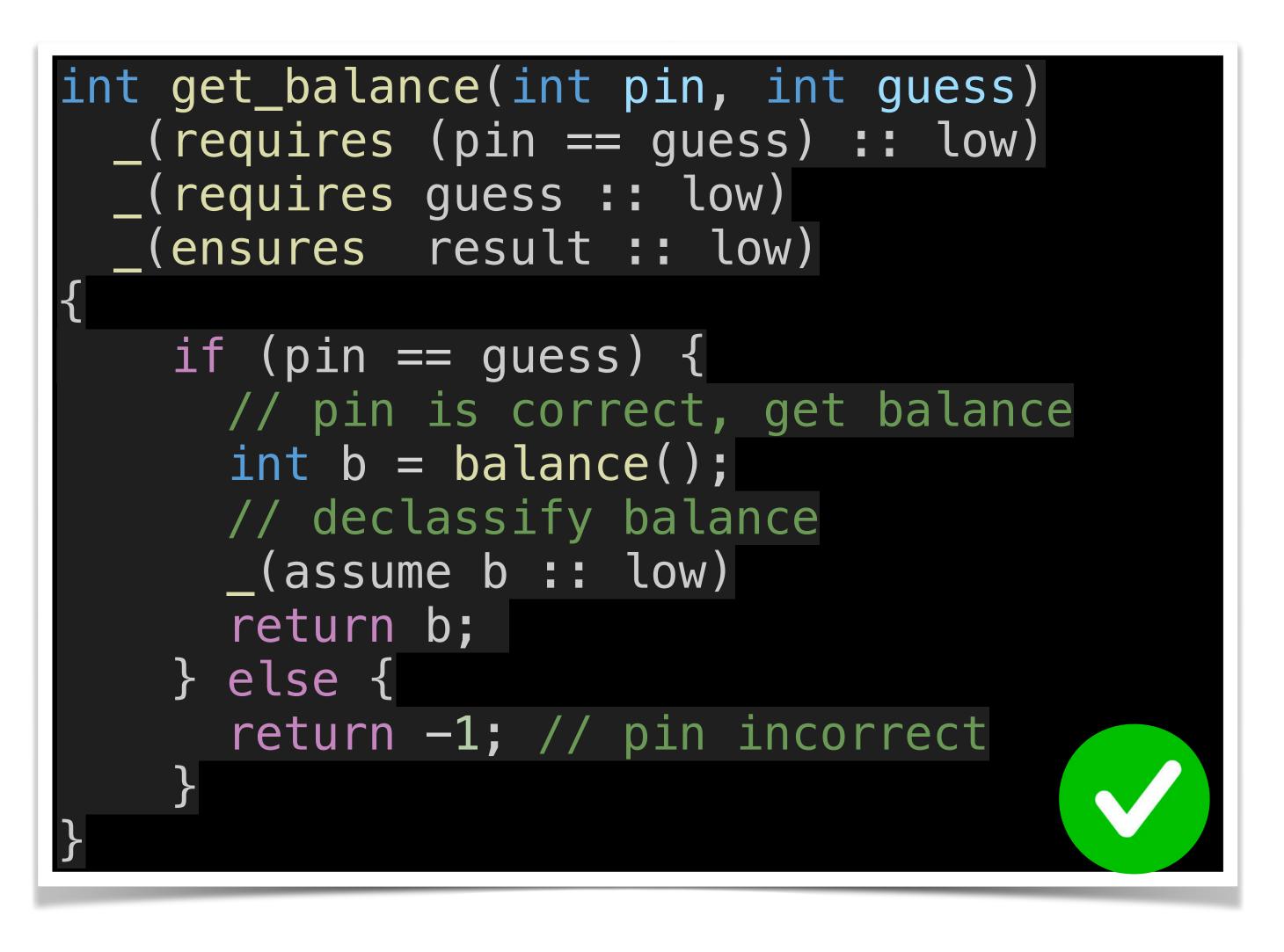


Declassification via Assume



$\vdash_{\ell} \{ \mathbf{emp} \} \mathbf{assume} \ \rho \ \{ \rho \}$

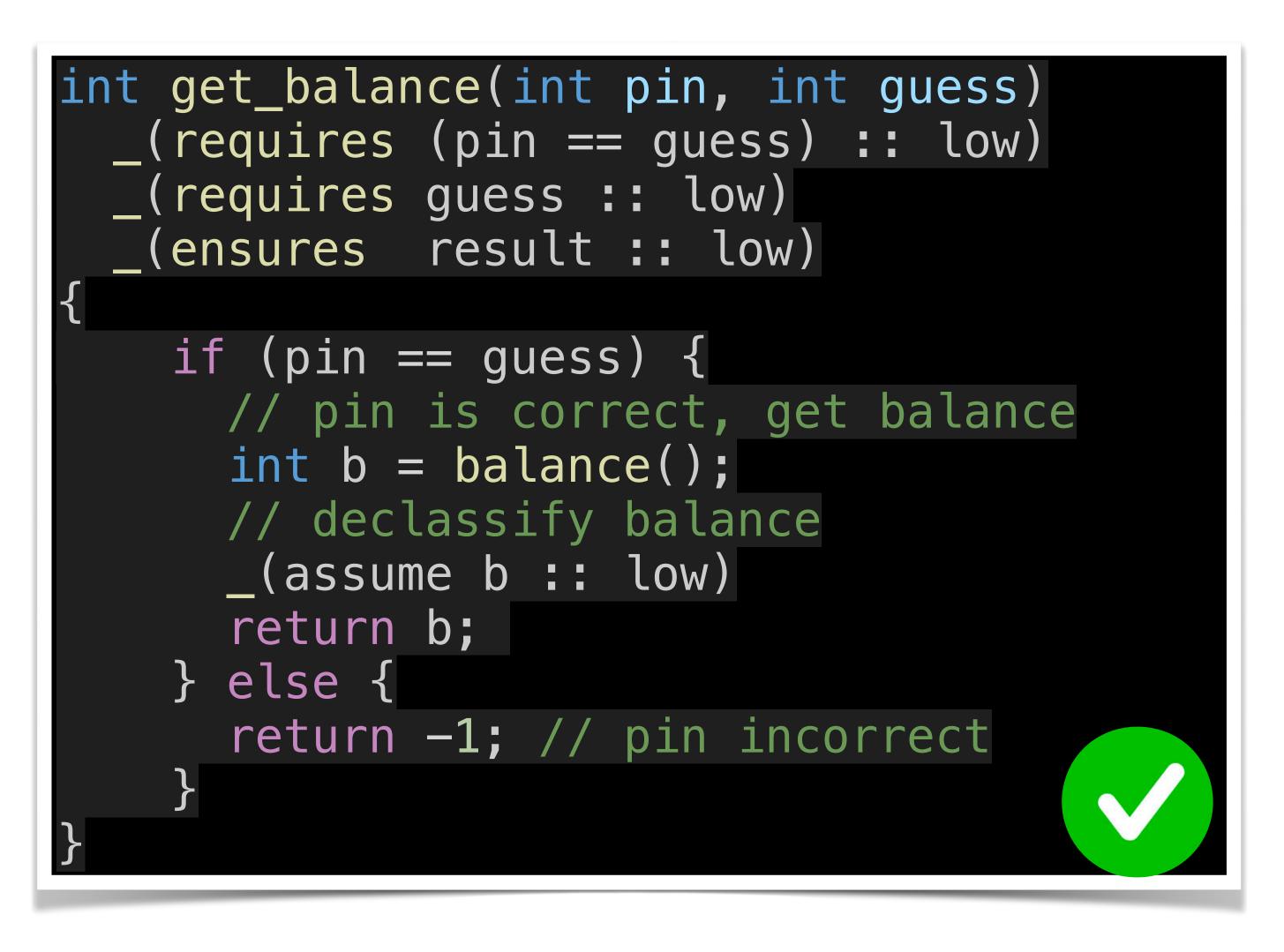
Declassification via Assume



This is magic

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Declassification via Assume



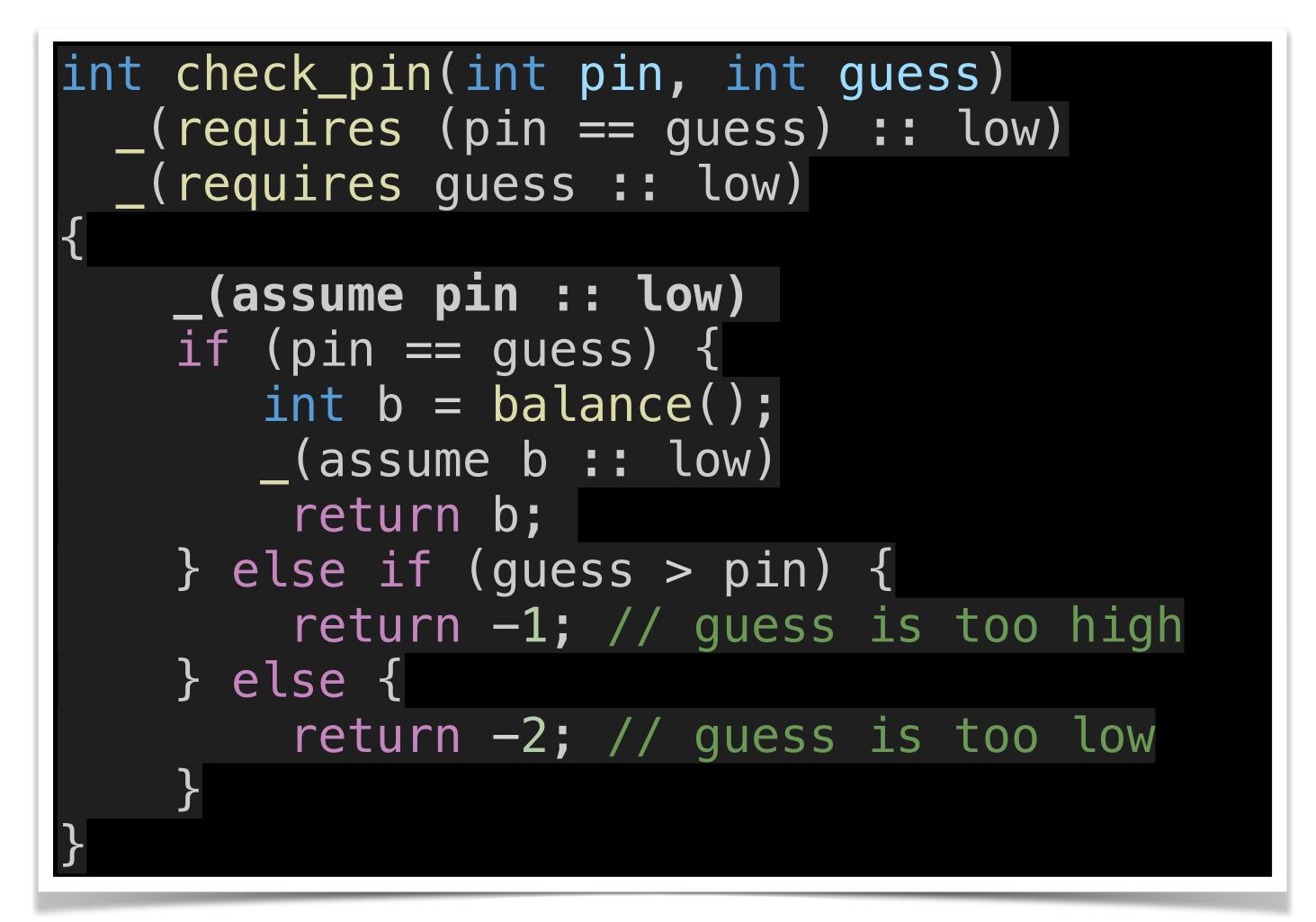
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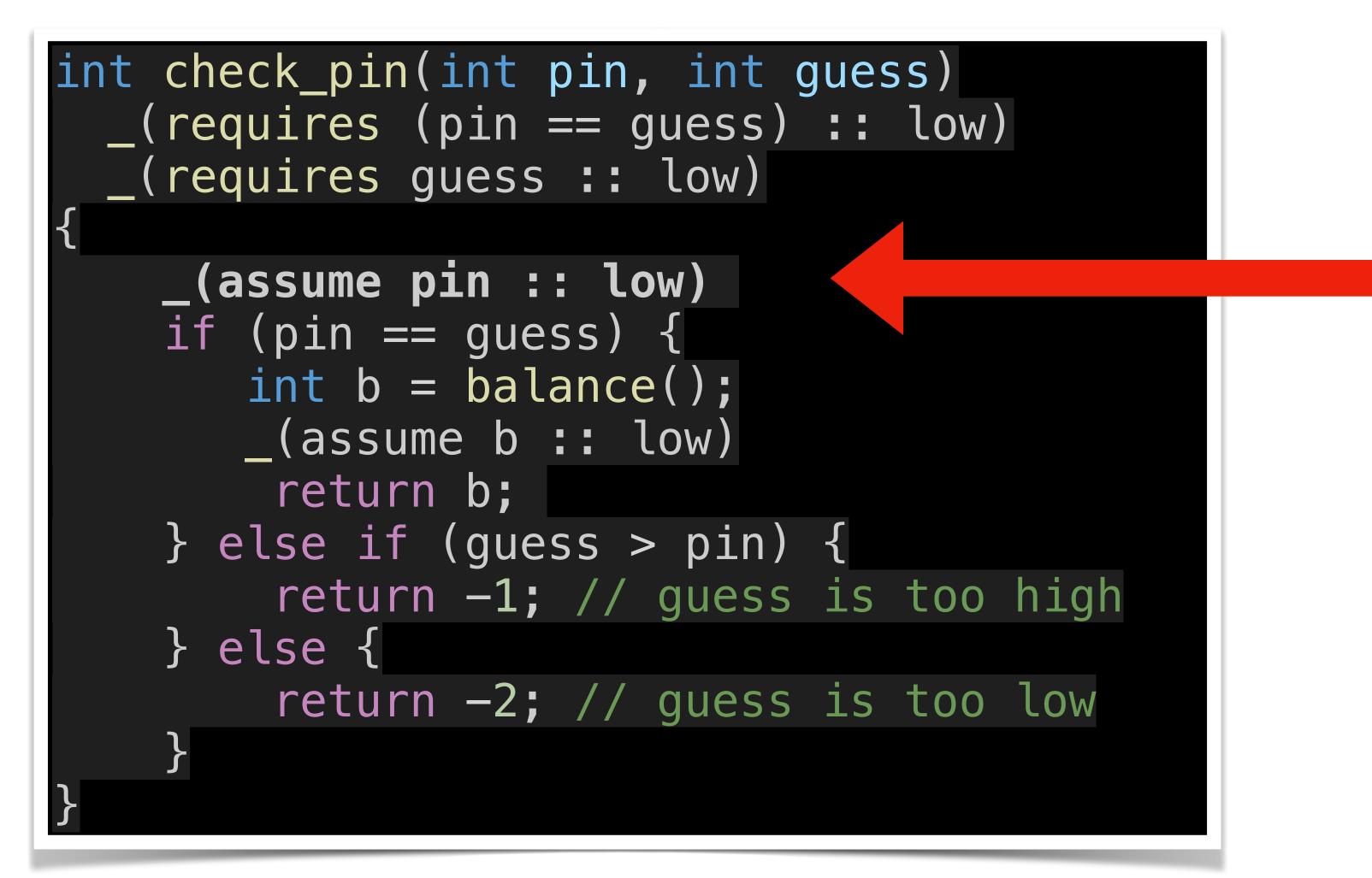
Contribution 1:

Semantic characterisation of soundness of this rule

Assume considered dangerous

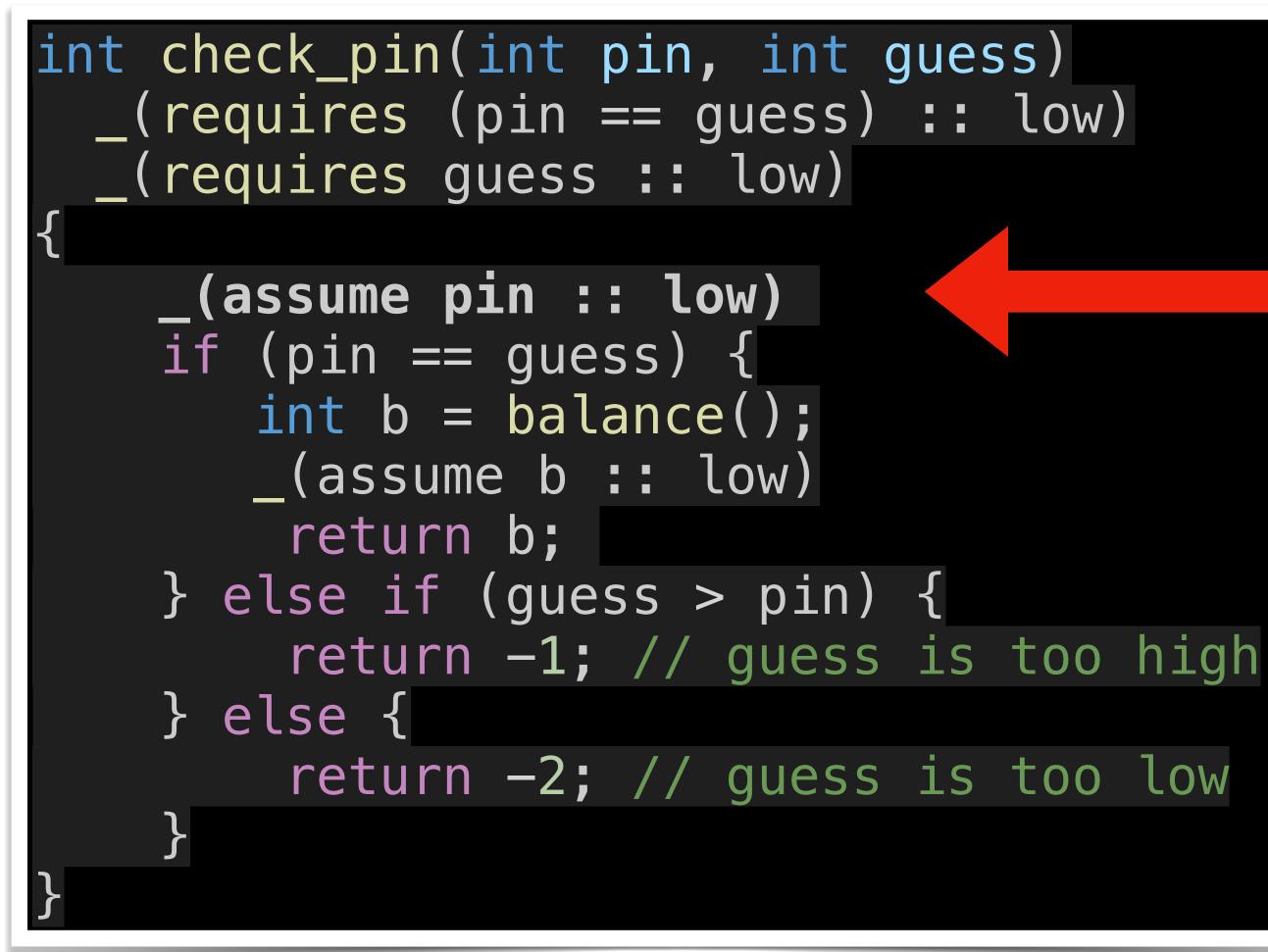


Assume considered dangerous



Insecure Declassification

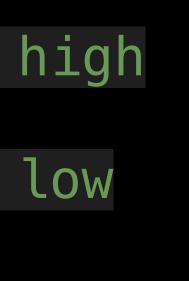
Assume considered dangerous





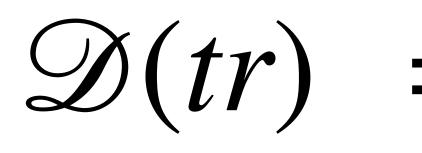


Justify declassifications against a declarative policy



Declarative Security Policies

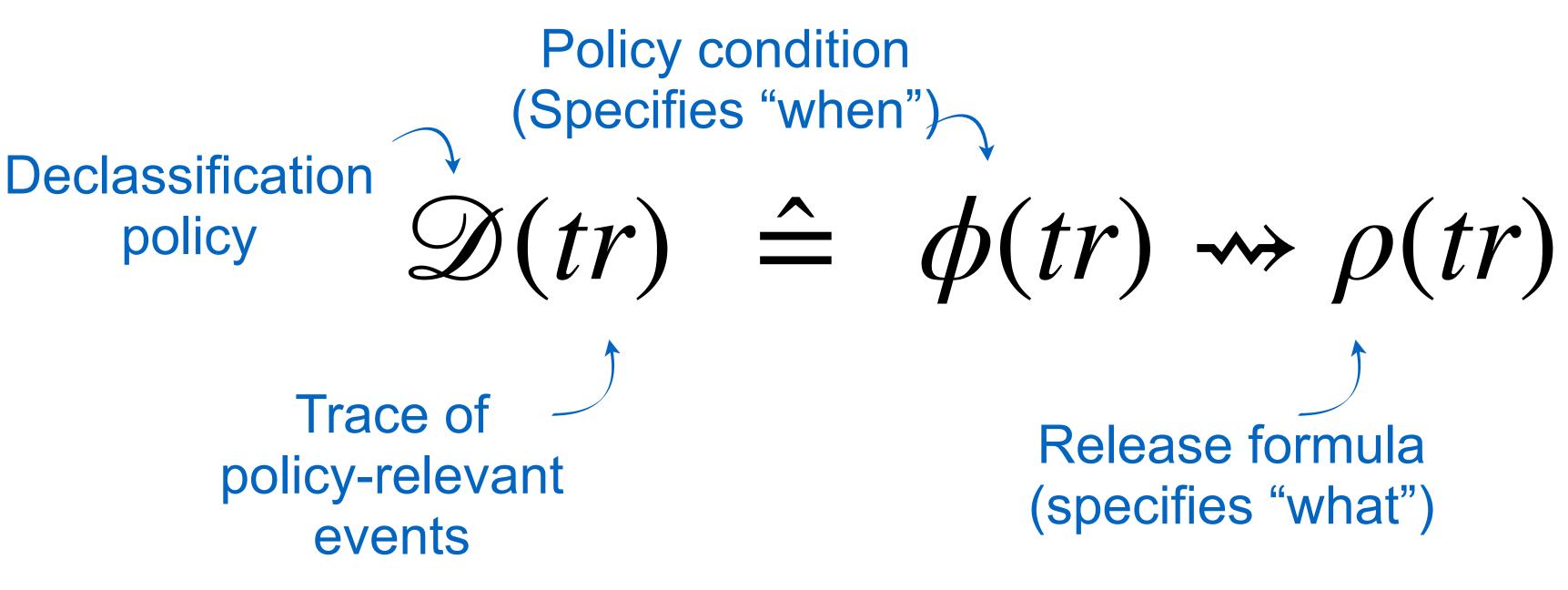
Specify when declassifications are allowed to occur and what information is allowed to be declassified when they do occur



 $\mathscr{D}(tr) \stackrel{\sim}{=} \phi(tr) \rightsquigarrow \rho(tr)$

Declarative Security Policies

Specify when declassifications are allowed to occur and what information is allowed to be declassified when they do occur



Release formula (specifies "what")

int check_pin(int pin, int guess) _(requires (pin == guess) :: low) _(requires guess :: low) _(assume pin :: low) if (pin == guess) { int b = balance(); _(assume b :: low) return b; } else if (guess > pin) { return -1; // too high } else { return -2; // too low

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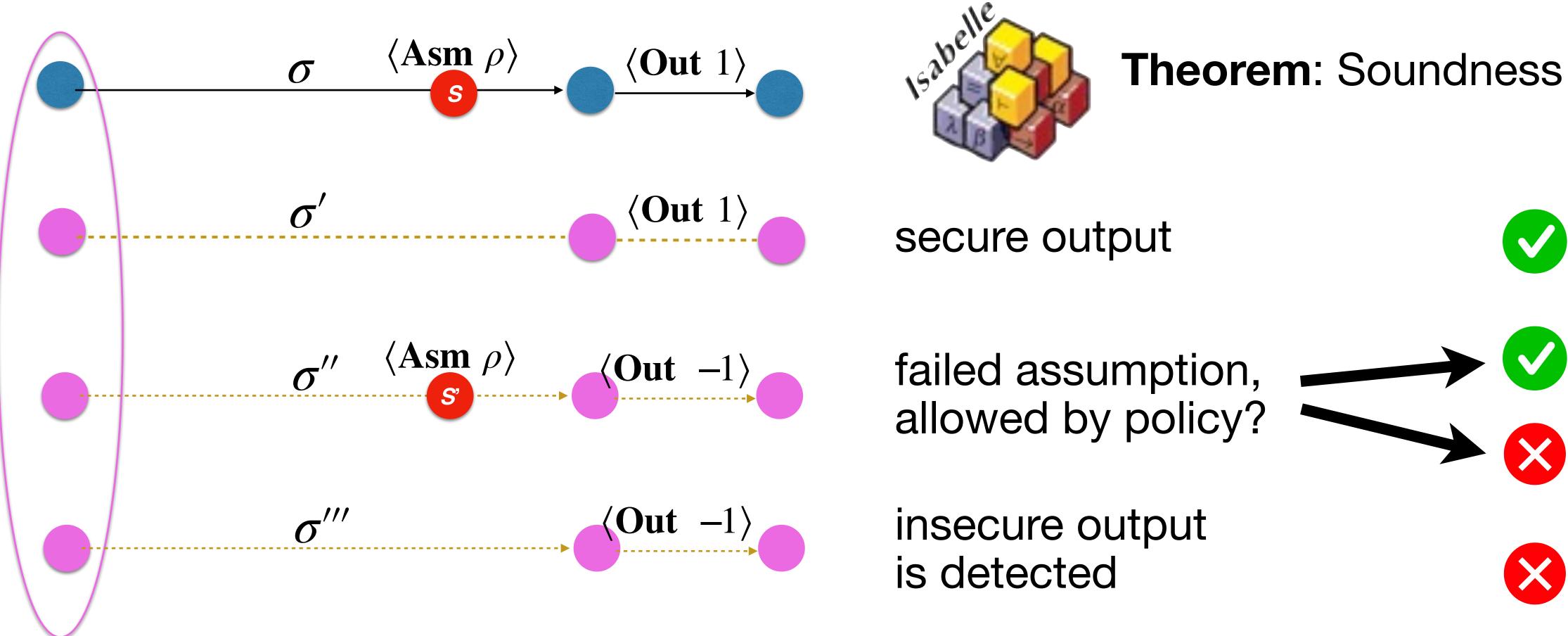
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What does Assume mean?

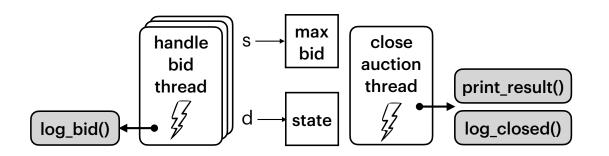
Policy-Agnostic Guarantee: Each step is allowed to leak some information, specifically whether an (earlier) assumption has been violated



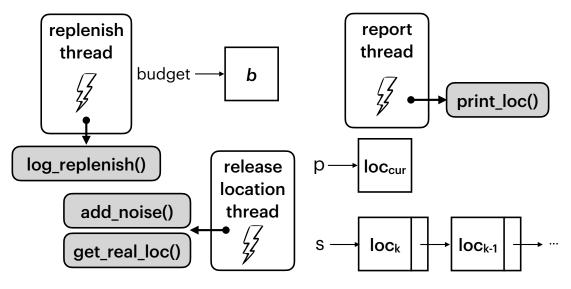
Case Studies



Sealed Bid Auction Server



Private Location Service





Private Learning

 $\mathscr{D}(tr) \stackrel{\sim}{=} \phi(tr) \rightsquigarrow \rho(tr)$

When: auction is finished

What: winning bid

When: sufficient privacy budget remains

What: a noisy location point (to ensure differential privacy)

https://verse.systems/blog/

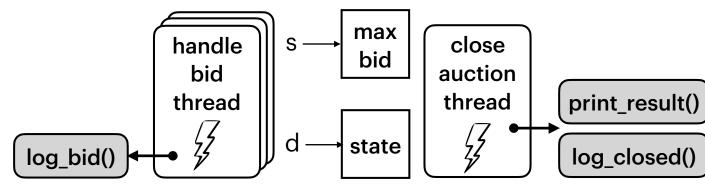
Thank You



https://covern.org/

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